

# The Linear Quadratic Regulator

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In these notes, we will derive the solution to the finite-horizon linear quadratic regulator (LQR) problem in several different ways. Fundamentally, LQR can be viewed as a large least-squares problem, but we are interested in the recursive solution because it can be efficiently computed (storage and computation scale linearly with the length of the time horizon).

## 1 The LQR problem

We consider the discrete-time finite-horizon version of the LQR problem. Consider the dynamical system with initial state  $x_0$  and

$$x_{t+1} = Ax_t + Bu_t \quad \text{for } t = 0, \dots, N-1 \quad (1)$$

The objective is to find a sequence of decisions  $u_0, \dots, u_{N-1}$  that minimizes the quadratic cost

$$J = \sum_{t=0}^{N-1} \underbrace{\begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix}}_{\text{stage cost}} + \underbrace{x_N^\top Q_f x_N}_{\text{terminal cost}} \quad (2)$$

The only assumptions we make are that  $\begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \succeq 0$ ,  $Q_f \succeq 0$ , and  $R \succ 0$ . These assumptions ensure that the cost will remain bounded. We first state result, and then we derive it in many ways.

**Theorem 1.** *The optimal decisions that solve the LQR problem are given by the state feedback policy  $u_t = K_t x_t$  for  $t = 0, \dots, N-1$ . We can compute the optimal policy recursively in an offline fashion by starting at  $t = N$  and working backwards to  $t = 0$ . The recursion is:*

$$P_N = Q_f \quad (3a)$$

$$P_t = A^\top P_{t+1} A + Q - (A^\top P_{t+1} B + S)(B^\top P_{t+1} B + R)^{-1}(B^\top P_{t+1} A + S^\top) \quad (3b)$$

$$K_t = -(B^\top P_{t+1} B + R)^{-1}(B^\top P_{t+1} A + S^\top) \quad (3c)$$

*The optimal cost starting from initial condition  $x_0$  is given by  $J_\star = x_0^\top P_0 x_0$ .*

**Note:** We can make the state and cost matrices time-varying if we like, i.e.  $A_t, B_t, Q_t, S_t, R_t$ . The solution is exactly analogous. We just have to make the recursion time-varying. So:

$$P_t = A_t^\top P_{t+1} A_t + Q_t - (A_t^\top P_{t+1} B_t + S_t)(B_t^\top P_{t+1} B_t + R_t)^{-1}(B_t^\top P_{t+1} A_t + S_t^\top)$$

$$K_t = -(B_t^\top P_{t+1} B_t + R_t)^{-1}(B_t^\top P_{t+1} A_t + S_t^\top)$$

In fact, we can even *make the sizes of all matrices time-varying!* For example, the state  $x_t$  and input  $u_t$  could have different sizes as  $t$  changes.

## 1.1 Solution via dynamic programming

Define the cost-to-go (optimal value function) for  $k = 0, \dots, N$  as

$$V_k(z) := \underset{u_k, \dots, u_{N-1}}{\text{minimize}} \quad \sum_{t=k}^{N-1} \begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix} + x_N^\top Q_f x_N$$

s.t.  $x_{t+1} = Ax_t + Bu_t$  for  $t = k, \dots, N-1$   
 $x_k = z$

Our ultimate goal is to find  $V_0(x_0)$ , but we will solve for all  $V_k$  for  $k = 0, \dots, N$ . By defining  $w := u_k$  and decomposing the value function by separating the first decision at time  $k$  from all subsequent decisions, we can show that the following recursive equation (the Bellman equation) holds:

$$V_k(z) = \min_w \left( \begin{bmatrix} z \\ w \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} z \\ w \end{bmatrix} + V_{k+1}(Az + Bw) \right) \quad \text{for } k = 0, \dots, N-1. \quad (4)$$

When  $k = N$ , we have  $V_N(z) = z^\top Q_f z$ . We can show by induction that  $V_k(z)$  is a positive semidefinite quadratic for all  $k \leq N$ . Suppose that  $V_t(z) = z^\top P_t z$  with  $P_t \succeq 0$  for  $t = k+1$ . We will prove that this holds for  $t = k$  as well. Substitute into Eq. (4) and obtain

$$V_k(z) = \min_w \left( \begin{bmatrix} z \\ w \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} z \\ w \end{bmatrix} + (Az + Bw)^\top P_{k+1} (Az + Bw) \right) \quad (5)$$

$$= \min_w \begin{bmatrix} z \\ w \end{bmatrix}^\top \begin{bmatrix} A^\top P_{k+1} A + Q & A^\top P_{k+1} B + S \\ B^\top P_{k+1} A + S^\top & B^\top P_{k+1} B + R \end{bmatrix} \begin{bmatrix} z \\ w \end{bmatrix} \quad (6)$$

This is a standard quadratic optimization problem. Due to our assumption that  $P_{k+1} \succeq 0$  and  $R \succ 0$ , the solution is

$$w^* = -(B^\top P_{k+1} B + R)^{-1} (B^\top P_{k+1} A + S^\top) z$$

$$V_k(z) = z^\top \left( A^\top P_{k+1} A + Q - (A^\top P_{k+1} B + S) (B^\top P_{k+1} B + R)^{-1} (B^\top P_{k+1} A + S^\top) \right) z$$

We deduce that  $V_k(z)$  is also quadratic, and  $P_k$  satisfies the recursion (3a)–(3b). Since  $w = u_k$  and  $z = x_k$ , we also find that the optimal policy is a state-feedback policy of the form  $u_t = K_t x_t$ , where  $K_t$  is given by (3c). The cost associated with using the optimal control policy starting from the state  $x_0$  is the cost to go  $V_0(x_0)$ , which is given by  $x_0^\top P_0 x_0$ .

**Note.** We assumed  $\begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \succeq 0$  and  $R \succ 0$ , so we can prove by induction that since  $P_N = Q_f \succeq 0$ , each  $V_t(z) = z^\top P_t z$  is the minimum of a positive definite quadratic function (5), and is therefore positive semidefinite, and we have  $P_t \succeq 0$  for all  $t$ .

The above dynamic programming approach works even when the system matrices are time-varying or even have different sizes as a function of time.

## 1.2 Solution via completing the square

Consider the cost we are trying to minimize:

$$J(x_0) = \sum_{t=0}^{N-1} \begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix} + x_N^\top Q_f x_N$$

Let's introduce a set of matrices  $P_0, P_1, \dots, P_N$  and include them into the sum as follows.

$$J(x_0) = x_0^\top P_0 x_0 + \sum_{t=0}^{N-1} \left( x_{t+1} P_{t+1} x_{t+1} - x_t^\top P_t x_t + \begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix} \right) + x_N^\top (Q_f - P_N) x_N.$$

Note that all the  $P_t$ 's cancel out, so the above expression is equal to  $J(x_0)$  *no matter what values we pick* for the  $P_t$ 's. Start by substituting  $x_{t+1} = Ax_t + Bu_t$  in the sum and it becomes

$$J(x_0) = x_0^\top P_0 x_0 + \sum_{t=0}^{N-1} \begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} A^\top P_{t+1} A - P_t + Q & A^\top P_{t+1} B + S \\ B^\top P_{t+1} A + S^\top & B^\top P_{t+1} B + R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix} + x_N^\top (Q_f - P_N) x_N.$$

Recall the completion of squares formula (LDU factorization):

$$\begin{bmatrix} x \\ u \end{bmatrix}^\top \begin{bmatrix} A & B \\ B^\top & C \end{bmatrix} \begin{bmatrix} x \\ u \end{bmatrix} = x^\top (A - BC^{-1}B^\top) x + (u - C^{-1}B^\top x)^\top C (u - C^{-1}B^\top x)$$

Applying this to the quadratic form in the sum, we obtain:

$$\begin{aligned} J(x_0) &= x_0^\top P_0 x_0 \\ &+ \sum_{t=0}^{N-1} x_t^\top \left( A^\top P_{t+1} A - P_t + Q - (A^\top P_{t+1} B + S)(B^\top P_{t+1} B + R)^{-1}(B^\top P_{t+1} A + S^\top) \right) x_t \\ &\quad + \sum_{t=0}^{N-1} (u_t - K_t x_t)^\top (B^\top P_{t+1} B + R) (u_t - K_t x_t) + x_N^\top (Q_f - P_N) x_N \end{aligned}$$

where we defined  $K_t$  as in (3c). Again, remember that this expression for  $J(x_0)$  *does not depend on the choice of the  $P_t$ 's*. So we can choose them however we like. In particular, if we choose  $P_t$  so that it satisfies (3a)–(3b), the sum simplifies greatly to

$$J(x_0) = x_0^\top P_0 x_0 + \sum_{t=0}^{N-1} (u_t - K_t x_t)^\top (B^\top P_{t+1} B + R) (u_t - K_t x_t). \quad (7)$$

We also have  $P_t \succeq 0$  for all  $t$  (see the note at the end of Section 1.1). Therefore each term in the sum is nonnegative. We can minimize  $J(x_0)$  by picking  $u_t = K_t x_t$ , which leaves us with the optimal cost  $J_\star = x_0^\top P_0 x_0$ .

**Note.** If we use a *suboptimal* policy  $\hat{K}_t$  instead of the optimal  $K_t$ , then the formula (7) reveals exactly the extra cost we will have to pay. In particular,

$$J_{\text{extra}} = \sum_{t=0}^{N-1} x_t^\top (\hat{K}_t - K_t)^\top (B^\top P_{t+1} B + R) (\hat{K}_t - K_t) x_t$$

### 1.3 Solution via block elimination

We will make use of *block variable elimination*. Here is a useful result that is easy to prove.

**Proposition 1** (block elimination). *Suppose we have linear equations of the form*

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} p \\ 0 \end{bmatrix},$$

where  $D$  is square and invertible. If we solve for  $y$  in the second equation and substitute the result into the first equation, we obtain

$$(A - BD^{-1}C)x = p \quad \text{and} \quad y = -D^{-1}Cx.$$

We will make use of this result throughout the following derivation.

Write out the objective and all constraints as a large optimization problem. Here, we treat both the states and inputs as variables, and we include the state dynamics as constraints.

$$\begin{aligned} & \underset{\substack{x_1, \dots, x_N, \\ u_0, \dots, u_{N-1}}}{\text{minimize}} && \sum_{t=0}^{N-1} \begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix} + x_N^\top Q_f x_N \\ & \text{s.t.} && x_{t+1} = Ax_t + Bu_t \quad \text{for } t = 0, \dots, N-1 \end{aligned}$$

Assign the Lagrange multiplier  $\lambda_{t+1}$  to the equality constraints for  $t = 0, \dots, N-1$ . The Lagrangian for the problem is therefore:

$$L(x, u, \lambda) = \frac{1}{2} \sum_{t=0}^{N-1} \begin{bmatrix} x_t \\ u_t \end{bmatrix}^\top \begin{bmatrix} Q & S \\ S^\top & R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix} + \frac{1}{2} x_N^\top Q_f x_N - \sum_{t=0}^{N-1} \lambda_{t+1}^\top (x_{t+1} - Ax_t - Bu_t)$$

The factors of  $\frac{1}{2}$  are there to make the algebra nicer. The KKT necessary conditions for optimality are  $\nabla_x L = 0$ ,  $\nabla_u L = 0$ , and  $\nabla_\lambda L = 0$ . Evaluating these gradients, we obtain the equations

$$\begin{aligned} Qx_t + Su_t + A^\top \lambda_{t+1} - \lambda_t &= 0 && \text{for } t = 0, \dots, N-1 \\ Q_f x_N - \lambda_N &= 0 \\ S^\top x_t + Ru_t + B^\top \lambda_{t+1} &= 0 && \text{for } t = 0, \dots, N-1 \\ Ax_t + Bu_t - x_{t+1} &= 0 && \text{for } t = 0, \dots, N-1 \end{aligned}$$

Merging these together as a single set of linear equations, we obtain:

$$\lambda_N = Q_f x_N \tag{8a}$$

$$\begin{bmatrix} \lambda_t \\ 0 \\ x_{t+1} \end{bmatrix} = \begin{bmatrix} Q & S & A^\top \\ S^\top & R & B^\top \\ A & B & 0 \end{bmatrix} \begin{bmatrix} x_t \\ u_t \\ \lambda_{t+1} \end{bmatrix} \quad \text{for } t = 0, \dots, N-1 \tag{8b}$$

We will prove by induction that  $\lambda_t = P_t x_t$  for all  $t$ . From (8a), the result holds for  $t = N$  with  $P_N = Q_f$ . Suppose it holds for  $t + 1$ . Substitute  $\lambda_{t+1} = P_{t+1} x_{t+1}$  into (8b) and obtain:

$$\begin{bmatrix} \lambda_t \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} Q & S & A^\top P_{t+1} \\ S^\top & R & B^\top P_{t+1} \\ A & B & -I \end{bmatrix} \begin{bmatrix} x_t \\ u_t \\ x_{t+1} \end{bmatrix} \tag{9}$$

Apply Proposition 1 to eliminate  $x_{t+1}$  from (9), which leads to:

$$\begin{bmatrix} \lambda_t \\ 0 \end{bmatrix} = \begin{bmatrix} A^\top P_{t+1} A + Q & A^\top P_{t+1} B + S \\ B^\top P_{t+1} A + S^\top & B^\top P_{t+1} B + R \end{bmatrix} \begin{bmatrix} x_t \\ u_t \end{bmatrix}$$

Apply Proposition 1 once more to eliminate  $u_t$ , which leads to:

$$\begin{aligned} \lambda_t &= \left( A^\top P_{t+1} A + Q - (A^\top P_{t+1} B + S)(B^\top P_{t+1} B + R)^{-1}(B^\top P_{t+1} A + S^\top) \right) x_t \\ u_t &= -(B^\top P_{t+1} B + R)^{-1}(B^\top P_{t+1} A + S^\top)x_t \end{aligned}$$

Therefore, we have  $\lambda_t = P_t x_t$ , which is what we wanted to prove, and the recursion for  $P_t$  and the expression for  $K_t$  are precisely the solution we previously found in Eq. (3).

**Alternative elimination ordering.** If we eliminate the variables in a different order, we get different (but equivalent) expressions for the  $P_t$  recursion and for  $K_t$ . Specifically, if we start from (9) but apply Proposition 1 to eliminate  $u_t$  first, we obtain:

$$\begin{aligned} \begin{bmatrix} \lambda_t \\ 0 \end{bmatrix} &= \begin{bmatrix} Q - SR^{-1}S^\top & A^\top P_{t+1} - SR^{-1}B^\top P_{t+1} \\ A - BR^{-1}S^\top & -I - BR^{-1}B^\top P_{t+1} \end{bmatrix} \begin{bmatrix} x_t \\ x_{t+1} \end{bmatrix} \\ u_t &= -R^{-1}(S^\top x_t + B^\top P_{t+1} x_{t+1}) \end{aligned}$$

To ease the notation, define:

$$E := A - BR^{-1}S^\top \quad G := BR^{-1}B^\top \quad \bar{Q} := Q - SR^{-1}S^\top$$

Based on our original problem assumptions, we have  $G \succeq 0$  and  $\bar{Q} \succeq 0$ . Using our new variable definitions, the equations simplify to:

$$\begin{aligned} \begin{bmatrix} \lambda_t \\ 0 \end{bmatrix} &= \begin{bmatrix} \bar{Q} & E^\top P_{t+1} \\ E & -(I + GP_{t+1}) \end{bmatrix} \begin{bmatrix} x_t \\ x_{t+1} \end{bmatrix} \\ u_t &= -R^{-1}(S^\top x_t + B^\top P_{t+1} x_{t+1}) \end{aligned}$$

Now apply Proposition 1 to eliminate  $x_{t+1}$  and obtain:

$$\begin{aligned} \lambda_t &= \left( \bar{Q} + E^\top P_{t+1} (I + GP_{t+1})^{-1} E \right) x_t \\ u_t &= -R^{-1}(S^\top + B^\top P_{t+1} (I + GP_{t+1})^{-1} E) x_t \\ x_{t+1} &= (I + GP_{t+1})^{-1} E x_t \end{aligned}$$

This yields new (but equivalent!) formulas for the optimal controller (3) and the optimal closed-loop matrix  $A + BK_t$ .

$$\boxed{\begin{aligned} P_N &= Q_f \\ P_t &= \bar{Q} + E^\top P_{t+1} (I + GP_{t+1})^{-1} E \\ K_t &= -R^{-1}(S^\top + B^\top P_{t+1} (I + GP_{t+1})^{-1} E) \\ A + BK_t &= (I + GP_{t+1})^{-1} E \end{aligned}} \tag{10}$$

## 1.4 Solution via adjoint equations

This approach is similar to the block elimination approach of Section 1.3. We start with (8):

$$\lambda_N = Q_f x_N \quad (11a)$$

$$\begin{bmatrix} \lambda_t \\ 0 \\ x_{t+1} \end{bmatrix} = \begin{bmatrix} Q & S & A^\top \\ S^\top & R & B^\top \\ A & B & 0 \end{bmatrix} \begin{bmatrix} x_t \\ u_t \\ \lambda_{t+1} \end{bmatrix} \quad \text{for } t = 0, \dots, N-1 \quad (11b)$$

Eliminate  $u_t$  right away using Proposition 1 and use the same new variables as in Section 1.3:

$$E := A - BR^{-1}S^\top \quad G := BR^{-1}B^\top \quad \bar{Q} := Q - SR^{-1}S^\top$$

This yields the so-called *adjoint equations*:

$$\lambda_N = Q_f x_N \quad (12a)$$

$$\begin{bmatrix} x_{t+1} \\ \lambda_t \end{bmatrix} = \begin{bmatrix} E & -G \\ \bar{Q} & E^\top \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_{t+1} \end{bmatrix} \quad (12b)$$

This is a difference equation with the state  $x_t$  equation evolving forward in time and co-state  $\lambda_t$  equation evolving backward in time. There is also a boundary condition that couples the variables at the terminal timestep. From here, we could prove  $\lambda_t = P_t x_t$  using induction as we did in Section 1.3. Another approach is to rearrange (12) so that both equations go forward in time, which yields

$$\lambda_N = Q_f x_N \quad (13a)$$

$$\begin{bmatrix} I & G \\ 0 & E^\top \end{bmatrix} \begin{bmatrix} x_{t+1} \\ \lambda_{t+1} \end{bmatrix} = \begin{bmatrix} E & 0 \\ -\bar{Q} & I \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t \end{bmatrix} \quad (13b)$$

If  $E$  is invertible, we can invert the matrix on the left-hand side and write the equations as

$$\lambda_N = Q_f x_N$$

$$\begin{bmatrix} x_{t+1} \\ \lambda_{t+1} \end{bmatrix} = \begin{bmatrix} E + GE^{-\top}\bar{Q} & -GE^{-\top} \\ -E^{-\top}\bar{Q} & E^{-\top} \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t \end{bmatrix}$$

The  $2 \times 2$  block matrix above is a *symplectic matrix* and has some useful properties, such as if  $\lambda$  is an eigenvalue, so is  $\lambda^{-1}$ . Such matrices play an important role in the study of Algebraic Riccati Equations. Consider a set of matrices  $P_0, P_1, \dots, P_N$  and write:

$$\begin{aligned} & \begin{bmatrix} x_{t+1} \\ \lambda_{t+1} - P_{t+1}x_{t+1} \end{bmatrix} \\ &= \begin{bmatrix} I & 0 \\ -P_{t+1} & I \end{bmatrix} \begin{bmatrix} x_{t+1} \\ \lambda_{t+1} \end{bmatrix} \\ &= \begin{bmatrix} I & 0 \\ -P_{t+1} & I \end{bmatrix} \begin{bmatrix} E + GE^{-\top}\bar{Q} & -GE^{-\top} \\ -E^{-\top}\bar{Q} & E^{-\top} \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t \end{bmatrix} \\ &= \begin{bmatrix} I & 0 \\ -P_{t+1} & I \end{bmatrix} \begin{bmatrix} E + GE^{-\top}\bar{Q} & -GE^{-\top} \\ -E^{-\top}\bar{Q} & E^{-\top} \end{bmatrix} \begin{bmatrix} I & 0 \\ P_t & I \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t - P_t x_t \end{bmatrix} \\ &= \begin{bmatrix} E + GE^{-\top}\bar{Q} - GE^{-\top}P_t & -GE^{-\top} \\ -P_{t+1}E - P_{t+1}GE^{-\top}\bar{Q} + P_{t+1}GE^{-\top}P_t + E^{-\top}P_t - E^{-\top}\bar{Q} & P_{t+1}GE^{-\top} + E^{-\top} \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t - P_t x_t \end{bmatrix} \end{aligned}$$

Note that this holds for *any* choice of the  $P_t$ , since we added and subtracted it without changing anything. Consider the (2, 1) block of the transition matrix:

$$\begin{aligned} & -P_{t+1}E - P_{t+1}GE^{-\top}\bar{Q} + P_{t+1}GE^{-\top}P_t + E^{-\top}P_t - E^{-\top}\bar{Q} \\ & = -P_{t+1}E + (P_{t+1}G + I)E^{-\top}(P_t - \bar{Q}) \end{aligned}$$

This can be made zero if we choose  $P_t = \bar{Q} + E^{\top}P_{t+1}(I + GP_{t+1})^{-1}E$ , which is precisely the alternative form for the solution we derived in (10). With this choice, our adjoint equations become:

$$\begin{bmatrix} x_{t+1} \\ \lambda_{t+1} - P_{t+1}x_{t+1} \end{bmatrix} = \begin{bmatrix} E + GE^{-\top}\bar{Q} - GE^{-\top}P_t & -GE^{-\top} \\ 0 & P_{t+1}GE^{-\top} + E^{-\top} \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t - P_t x_t \end{bmatrix}$$

Substituting for  $P_t$  and simplifying, we obtain:

$$\begin{bmatrix} x_{t+1} \\ \lambda_{t+1} - P_{t+1}x_{t+1} \end{bmatrix} = \begin{bmatrix} (I + GP_{t+1})^{-1}E & -GE^{-\top} \\ 0 & (I + P_{t+1}G)E^{-\top} \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t - P_t x_t \end{bmatrix}$$

Now recall from (10) that  $A + BK_t = (I + GP_{t+1})^{-1}E$ , so we have:

$$\begin{bmatrix} x_{t+1} \\ \lambda_{t+1} - P_{t+1}x_{t+1} \end{bmatrix} = \begin{bmatrix} A + BK_t & -GE^{-\top} \\ 0 & (A + BK_t)^{-\top} \end{bmatrix} \begin{bmatrix} x_t \\ \lambda_t - P_t x_t \end{bmatrix}$$

From here, we easily see that if  $\lambda_{t+1} = P_{t+1}x_{t+1}$ , then we must also have  $\lambda_t = P_t x_t$  and this completes the proof. The equations also simplify to  $x_{t+1} = (A + BK_t)x_t$ , which are the closed-loop equations we expected to see.

**Infinte-horizon LQR.** This formulation using the adjoint equation is particularly useful when solving the infinite-horizon LQR problem. In the infinite-horizon setting, we have  $P_t = P_{t+1} = P$ , so the transformation of the symplectic matrix preserves eigenvalues, and we have:

$$\begin{bmatrix} I & 0 \\ -P & I \end{bmatrix} \underbrace{\begin{bmatrix} E + GE^{-\top}\bar{Q} & -GE^{-\top} \\ -E^{-\top}\bar{Q} & E^{-\top} \end{bmatrix}}_M \begin{bmatrix} I & 0 \\ P & I \end{bmatrix} = \begin{bmatrix} A + BK & -GE^{-\top} \\ 0 & (A + BK)^{-\top} \end{bmatrix}. \quad (14)$$

This observation is the key to solving the Discrete Algebraic Riccati Equation (DARE): eigenvalues of the symplectic matrix  $M$  are the eigenvalues of the LQR-optimal closed-loop map (stable) and their conjugate inverses (unstable). Multiply (14) by  $\begin{bmatrix} I & 0 \\ P & I \end{bmatrix} (\dots) \begin{bmatrix} I \\ 0 \end{bmatrix}$  and obtain

$$M \begin{bmatrix} I \\ P \end{bmatrix} = \begin{bmatrix} I \\ P \end{bmatrix} (A + BK). \quad (15)$$

The stable eigenvalues of  $M$  are the eigenvalues of  $(A + BK)$ . So if we diagonalize  $M$  and collect all stable eigenvalues in the diagonal matrix  $\Lambda$ , we can write the eigenvalue decomposition

$$M \begin{bmatrix} V_1 \\ V_2 \end{bmatrix} = \begin{bmatrix} V_1 \\ V_2 \end{bmatrix} \Lambda.$$

Under suitable assumptions,  $V_1$  will be invertible. Multiply on the right by  $V_1^{-1}$  and obtain

$$M \begin{bmatrix} I \\ V_2 V_1^{-1} \end{bmatrix} = \begin{bmatrix} I \\ V_2 V_1^{-1} \end{bmatrix} (V_1 \Lambda V_1^{-1}).$$

Note the similarity with (15). It takes some work to prove the details, but it turns out that  $P = V_2 V_1^{-1}$  is the (unique) stabilizing solution to the DARE, and  $V_1 \Lambda V_1^{-1} = A + BK$  is the LQR-optimal closed-loop map.